

Abelian Groups III

Today, we are going to set up the machinery we will use to prove:

Theorem. *If S is subgroup of \mathbb{Z}^n , then we can choose a new basis $\{b_1, \dots, b_n\}$ of \mathbb{Z}^n and a new generating set $\{t_1, \dots, t_\ell\}$ ($\ell \leq n$) for S such that for each $i = 1, \dots, \ell$, t_i is an integer multiple of b_i , i.e., $t_i = d_i b_i$ for some $d_i \in \mathbb{Z}$.*

This theorem ought to have a distinguished name—maybe, “Fundamental Theorem of Finitely-Generated Abelian Groups.” But there are numerous ways of stating it and its immediate consequences, and maybe this stands in the way of a standard name. The theorem is closely related to “Smith Normal Form,” but I will not work with this concept in the present lecture.¹

As I said at the end of the last lecture, the theorem has remarkable consequences. First, it means that every subgroup of a free abelian group is free, for the t_i form a basis for S . (A non-trivial \mathbb{Z} -linear relation among the t_i would produce a non-trivial relation among the b_i .) Second, the theorem implies that every finitely generated abelian group is a direct sum of cyclic subgroups. For suppose G is any finitely-generated abelian group. Then there is a surjection $\mathbb{Z}^n \rightarrow G$. Let S be the kernel of this map, and choose $\{b_1, \dots, b_n\}$ and $\{t_1, \dots, t_\ell\}$ as in the theorem. Then

$$S = \bigoplus_{i=1}^{\ell} \mathbb{Z}t_i,$$

and it follows (as we shall show in more detail later) that

$$G \cong \mathbb{Z}^n / S \cong \mathbb{Z}/d_1\mathbb{Z} \oplus \mathbb{Z}/d_2\mathbb{Z} \oplus \cdots \oplus \mathbb{Z}/d_\ell\mathbb{Z} \oplus \mathbb{Z}^{n-\ell}.$$

Observe that some of the d_i may be equal to 1, in which case the corresponding summand is $\{0\}$. Finally, let me point out that the proof of the theorem is constructive, allowing us to compute the basis $\{b_1, \dots, b_n\}$ and the d_i from any list of generators for S . The procedure that I describe below allows some choices to the user, but it can be turned into a computer program. Finding fast implementations is an area of current research.

Using matrices to describe morphisms between free R -modules.

Let R be a ring with multiplicative identity. We have shown that R^n is the free R -module on the standard basis $\{e_1, \dots, e_n\} \subset R^n$, where

$$e_{ij} = (e_i)_j = \pi_j(e_i) = \begin{cases} 1_R, & \text{if } i = j; \\ 0, & \text{otherwise.} \end{cases}$$

We are going to introduce a way of displaying the data required to describe elements and morphisms of free R -modules relative to a basis. It’s much like the notation introduced earlier for vector spaces,

¹ You might look up “Smith Normal Form” on the web. You can even read Smith’s original paper:

<http://www.jstor.org/stable/108738>.

but we switch the roles of rows and columns. This seems to be done mainly to remain consistent with the usual matrix multiplication conventions when R is not commutative.

Let $B = \{b_1, \dots, b_m\}$ be a basis for R^m . Each element $x \in R^m$ has a unique expression as $r_1 b_1 + \dots + r_m b_m$, with $r_i \in R$. The row vector with the r_i as entries is denoted

$$(x; B) := (r_1 \ \cdots \ r_m).$$

Suppose R^n has basis $C = \{c_1, \dots, c_n\}$ and $\phi: R^m \rightarrow R^n$. Let $(\phi; BC)$ denote the matrix whose m rows are the $(\phi(b_1); C), \dots, (\phi(b_m); C)$:

$$(\phi; BC) := \begin{pmatrix} (\phi(b_1); C) \\ \vdots \\ (\phi(b_m); C) \end{pmatrix}.$$

In other words, the entries of $(\phi; BC)$ are—in the i^{th} row and j^{th} column—the elements $s_{ij} \in R$ defined by

$$\phi(b_i) = s_{i1}c_1 + \dots + s_{in}c_n.$$

This results in conventions that are compatible with non-commutative rings. Indeed, if $x = r_1 b_1 + \dots + r_m b_m$, then

$$\begin{aligned} \phi(x) &= \phi(r_1 b_1 + \dots + r_m b_m) \\ &= r_1 \phi(b_1) + \dots + r_m \phi(b_m) \\ &= r_1 (s_{11}c_1 + \dots + s_{1n}c_n) + \dots + r_m (s_{m1}c_1 + \dots + s_{mn}c_n) \\ &= (r_1 s_{11} + r_2 s_{21} + \dots + r_m s_{m1})c_1 + \dots + (r_1 s_{1n} + r_2 s_{2n} + \dots + r_m s_{mn})c_n. \end{aligned}$$

This is consistent with matrix multiplication:

$$\begin{aligned} (x; B)(\phi; BC) &= (r_1 \ \cdots \ r_m) \begin{pmatrix} s_{11} & \cdots & s_{1n} \\ \vdots & \vdots & \vdots \\ s_{m1} & \cdots & s_{mn} \end{pmatrix} \\ &= (r_1 s_{11} + r_2 s_{21} + \dots + r_m s_{m1} \quad \cdots \quad r_1 s_{1n} + r_2 s_{2n} + \dots + r_m s_{mn}) \\ &= (\phi(x); C). \end{aligned}$$

Exercise. Suppose R is not commutative. What would go wrong in attempting to represent morphisms of free R -modules by matrix multiplication if we chose to represent elements by columns and the application of a morphism by matrix multiplication with the matrix on the left.

Suppose $\phi: R^m \rightarrow R^n$ is as above, R^p has basis D and $\psi: R^n \rightarrow R^p$. Then we have the following formulae:

$$\begin{aligned} (x; B)(\phi; BC)(\psi; CD) &= (\psi(\phi(x)); D); \\ (\phi; BC)(\psi; CD) &= (\psi\phi; BD). \end{aligned}$$

Just as previously described for vector spaces, if B and C are two bases for R^n , the matrix $(\text{id}_{R^n}; BC)$ effects a change of base:

$$(x; B)(\text{id}_{R^n}; BC) = (x; C).$$

The i^{th} row of $(\text{id}_{R^n}; BC)$ is the n -tuple of coefficients required to write b_i as an R -linear combination of the c_1, \dots, c_n . This is an $n \times n$ matrix with right inverse $(\text{id}_{R^n}; CB)$. In case R is commutative, these matrices are inverses. In the non-commutative case, there are complications: a matrix may have a right (resp., left) inverse that is not a left (resp., right) inverse; see Jacobson, *Basic Algebra I*, page 97, Exercise 2. Even worse, in the non-commutative case, it is possible for $R^m \cong R^n$ with $m \neq n$; see Jacobson, page 171.

Diagonalizing integer matrices

Suppose $A = \{a_{ij}\}$ is an $m \times n$ integer matrix.

How to add a multiple of one row (column) of A to another row (column). Let $T_{ij}^n(k)$, $i \neq j$, $i, j \in \{1, \dots, n\}$, be the $n \times n$ matrix with k in the $(i, j)^{th}$ position and 1s on the diagonal. The inverse of $T_{ij}^n(k)$ is $T_{ij}^n(-k)$. Note that $T_{ij}^n(k)A$ has the same rows as A , except for the i^{th} row, which is the i^{th} row of A plus k times the j^{th} row of A . Similarly, $AT_{ij}^n(k)$ has the same columns as A , except for the j^{th} column, which is the j^{th} column of A plus k times the i^{th} column of A .

How to switch two rows (columns) of A . Let P_{ij}^n be the $n \times n$ identity matrix with rows i and j switched. The same matrix arises by switching columns i and j . P_{ij}^n is its own inverse. Notice that $P_{ij}^m A$ is the same as A , but with rows i and j switched. AP_{ij}^n is the same as A , but with columns i and j switched.

How to multiply a row (column) by -1 . Let U_i^n be the diagonal matrix in which the ii^{th} entry is -1 and all others are 1. U_i^n is its own inverse. $U_i^m A$ is the same as A , but with its i^{th} row multiplied by -1 . AU_j^n is the same as A , but with its j^{th} column multiplied by -1 .

Column Step. By repeatedly multiplying A on the left by various matrices of the form $T_{ij}^m(k)$ and P_{ij}^m , we can bring it to the form

$$A' = \begin{pmatrix} a'_{11} & a'_{12} & \cdots & a'_{1n} \\ 0 & a'_{22} & \cdots & a'_{2n} \\ \vdots & \vdots & \vdots & \vdots \\ 0 & a'_{m2} & \cdots & a'_{mn} \end{pmatrix},$$

where the first column has zero in every place but the first and a'_{11} is the greatest common divisor of the entries in the first column of A . The reason that we can do this is essentially that the *GCD* of any finite set of integers can be expressed as a \mathbb{Z} -linear combination of them. I will not specify the precise steps used to select the matrices to multiply by. One might, for example, subtract (or add) a row with the smallest (in absolute value) initial entry from the other rows, and do this over and over until only one row with a non-zero initial entry remained. In examples that are done by hand, it may be more convenient or faster to use some other method. The specific numbers themselves may suggest a strategy.

Row Step. By repeatedly multiplying A on the right by various matrices of the form $T_{ij}^n(k)$ and P_{ij}^n , we can bring it to the form

$$A'' = \begin{pmatrix} a''_{11} & 0 & \cdots & 0 \\ a''_{21} & a''_{22} & \cdots & a''_{2n} \\ \vdots & \vdots & \vdots & \vdots \\ a''_{m1} & a''_{m2} & \cdots & a''_{mn} \end{pmatrix},$$

where the first row has zero in every place but the first and a''_{11} is the greatest common divisor of the entries in the first row of A .

We can perform these steps without multiplying rows or columns by -1 , but it may be useful or convenient to make some entries positive, and this can clearly be done by using U_i^n .

Barring the possibility that some $a_{11}^{(s)}$ is zero, $|a_{11}^{(s+1)}| < |a_{11}^{(s)}|$. Thus, by repeatedly performing row steps and column steps, we may bring A to the form

$$B = \begin{pmatrix} b_{11} & 0 & \cdots & 0 \\ 0 & b_{22} & \cdots & b_{2n} \\ \vdots & \vdots & \vdots & \vdots \\ 0 & b_{m2} & \cdots & b_{mn} \end{pmatrix},$$

where b_{11} is the greatest common divisor of all the elements in the first column and first row of A . If some $a_{11}^{(s)}$ is zero, we may either make it non-zero by exchanging t

We an then apply the same process to the submatrix $\begin{pmatrix} b_{22} & \cdots & b_{2n} \\ \vdots & \vdots & \vdots \\ b_{m2} & \cdots & b_{mn} \end{pmatrix}$. We can do this “in place” by multiplying B on the right and left by $T_{ij}^n(k)$ and P_{ij}^n with i and j never equal to 1. We get a matrix of the form:

$$B' = \begin{pmatrix} b_{11} & 0 & 0 & \cdots & 0 \\ 0 & c_{22} & 0 & \cdots & 0 \\ 0 & 0 & c_{33} & \cdots & c_{3n} \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & c_{m3} & \cdots & c_{mn} \end{pmatrix}.$$

Continuing, we get a diagonal matrix D and a relationship:

$$PAQ = D,$$

where P is $m \times m$, Q is $n \times n$ and each is a product of matrices of the form T_{ij} and P_{ij} and therefore is invertible.

Closing remarks. We have not required any conditions on the diagonal entries in PAQ . If the entries are such that $d_i | d_{i+1}$ for $i = 1, \dots, \ell$, $\ell \leq m$, and $d_i = 0$ for $i > \ell$, then we say that the matrix is in Smith Normal Form. We can assure that PAQ winds up in this form by a slightly more complex algorithm than the one we described, but we do not need this detail to complete our analysis of finitely-generated abelian groups.